

# On Hypercontractivity and a Data Processing Inequality

Venkat Anantharam  
EECS Department,  
University of California,  
Berkeley, CA, USA  
ananth@eecs.berkeley.edu

Amin Gohari  
ISSL Lab, EE Department,  
Sharif University of Technology,  
Tehran, Iran  
aminzadeh@sharif.edu

Sudeep Kamath  
ECE Department,  
University of California,  
San Diego, CA, USA  
sukamath@ucsd.edu

Chandra Nair  
IE Department,  
The Chinese University  
of Hong Kong  
chandra@ie.cuhk.edu.hk

**Abstract**—In this paper we provide the correct tight constant to a data-processing inequality claimed by Erkip and Cover. The correct constant turns out to be a particular hypercontractivity parameter of  $(X, Y)$ , rather than their squared maximal correlation. We also provide alternate geometric characterizations for both maximal correlation as well as the hypercontractivity parameter that characterizes the data-processing inequality.

## I. INTRODUCTION

Given a pair of random variables<sup>1</sup>  $(X, Y)$ , the *data-processing inequality* states that whenever  $U \rightarrow X \rightarrow Y$  form a Markov chain, we have

$$I(U; Y) \leq I(U; X).$$

A natural question to ask is the following: what is the smallest  $r$  such that the inequality

$$I(U; Y) \leq rI(U; X)$$

holds for every  $U$  whenever  $U \rightarrow X \rightarrow Y$  is Markov. Erkip and Cover [3] claimed that the smallest possible  $r$  is  $\rho_m^2(X; Y)$ , the squared *maximal correlation* between  $X$  and  $Y$ . We show that this result is incorrect and we establish that the right constant is related to a particular hypercontractivity parameter of  $(X, Y)$ .

### A. Definitions and preliminaries

**Definition 1.** For any real-valued random variable  $W$  with finite support, and any real number  $p \geq 1$ , define  $\|W\|_p := (\mathbb{E}|W|^p)^{\frac{1}{p}}$ .

**Definition 2.** Given random variables  $X$  and  $Y$ , the Hirschfeld-Gebelein-Rényi maximal correlation of  $(X, Y)$  is defined as follows:

$$\rho_m(X; Y) := \sup \mathbb{E}[f(X)g(Y)], \quad (1)$$

where the supremum is taken over all functions  $f, g$  such that

$$\mathbb{E} f(X) = \mathbb{E} g(Y) = 0, \text{ and } \mathbb{E} f^2(X), \mathbb{E} g^2(Y) \leq 1.$$

<sup>1</sup>Throughout this paper, random variables  $(X, Y)$  take values in  $\mathcal{X} \times \mathcal{Y}$  with  $|\mathcal{X}|, |\mathcal{Y}| < \infty$ . Further we assume that  $\mathbb{P}(X = x) > 0 \forall x \in \mathcal{X}$  and  $\mathbb{P}(Y = y) > 0 \forall y \in \mathcal{Y}$ .

**Definition 3.** A pair of random variables  $(X, Y)$  is said to be  $(p, q)$ -hypercontractive for  $1 \leq q \leq p < \infty$  if the inequality

$$\| \mathbb{E}(g(Y)|X) \|_p \leq \|g(Y)\|_q$$

holds for all functions  $g(Y)$ .

**Definition 4.** For  $p \geq 1$  define  $q_p^*(X; Y)$  as

$$\inf\{q : q \geq 1, (X, Y) \text{ is } (p, q) \text{ - hypercontractive}\}.$$

For  $p \geq 1$ , we define the following quantity:

$$r_p(X; Y) := \frac{q_p^*(X; Y)}{p}.$$

An important property of the hypercontractivity parameter  $r_p(X; Y)$  is the so-called *tensorization property*. It is known [1] that if  $(X_1, Y_1)$  is independent of  $(X_2, Y_2)$ , then

$$r_p(X_1, X_2; Y_1, Y_2) = \max\{r_p(X_1; Y_1), r_p(X_2; Y_2)\}.$$

The following theorem summarizes some results about the quantity  $r_p(X; Y)$ .

**Theorem 1** ([1] Theorem 3a,3b). *The following statements hold:*

- (i)  $r_p(X; Y)$  is non-increasing in  $p$ .
- (ii)  $r_p(X; Y) \geq \rho_m^2(X; Y) + \frac{1 - \rho_m^2(X; Y)}{p}$  for all  $p \geq 1$ .

Denote

$$r_\infty(X; Y) := \inf_{p \geq 1} r_p(X; Y). \quad (2)$$

**Remark 1.** It is clear from Theorem 1 that

$$r_\infty(X; Y) = \lim_{p \rightarrow \infty} r_p(X; Y) \geq \rho_m^2(X; Y). \quad (3)$$

Let  $\nu_X(x)$  and  $\mu_X(x)$  be probability distributions on the same finite set. We use  $D_{KL}(\nu_X \parallel \mu_X)$  to denote the relative entropy or the Kullback-Liebler divergence between  $\nu_X(x)$  and  $\mu_X(x)$ , i.e.

$$D_{KL}(\nu_X \parallel \mu_X) := \sum_x \nu_X(x) \log \frac{\nu_X(x)}{\mu_X(x)}.$$

Given a pair of random variables  $(X, Y) \sim \mu(x, y)$  where  $\mu(x, y)$  denotes their probability mass function, let  $\mu(y|x)$  be the channel from  $X$  to  $Y$  induced by  $\mu(x, y)$ . We consider  $X$  to be the input and  $Y$  to be the output of the channel. For any input  $\nu_X(x)$ , let  $\nu_Y^\mu(y) = \sum_x \nu_X(x)\mu(y|x)$  denote the induced output distribution by the channel  $\mu(y|x)$  when the input distribution is  $\nu_X$ . Define

$$d_*(X; Y) := \sup_{\nu_X \neq \mu_X} \frac{D_{KL}(\nu_Y^\mu \parallel \mu_Y)}{D_{KL}(\nu_X \parallel \mu_X)}. \quad (4)$$

Finally define the main quantity of interest:

$$m_*(X; Y) := \sup_{U: U-X-Y, I(U; X) > 0} \frac{I(U; Y)}{I(U; X)}. \quad (5)$$

Here again we think of  $X$  as the input and  $Y$  as the output of the channel characterized by  $\mu(y|x)$ .

**Remark 2.** *Witsenhausen and Wyner [13] consider the trade-off between possible values of  $I(U; X)$  and  $I(U; Y)$ . The tensorization property of  $m_*(X; Y)$  can be inferred from their results.*

### B. Summary of results

We provide a proof of the following equivalence result:

**Theorem 2.** *Given a pair of random variables  $(X, Y)$  and the quantities  $r_\infty(X; Y)$ ,  $d_*(X; Y)$ , and  $m_*(X; Y)$  as defined in the previous section, we have*

$$r_\infty(X; Y) = d_*(X; Y) = m_*(X; Y).$$

**Remark 3.** *The equality  $r_\infty(X; Y) = d_*(X; Y)$  was established in [1, Theorem 5a], and only the second equality  $d_*(X; Y) = m_*(X; Y)$  is new here. However we will prove a three way equivalence in this paper as opposed to only establishing  $d_*(X; Y) = m_*(X; Y)$ .*

**Remark 4.** *In [3, Theorem 8] it was claimed that the following inequality holds:*

$$I(U; Y) \leq \rho_m^2(X; Y)I(U; X), \quad \forall U - X - Y.$$

*It turns out that this inequality is incorrect. Indeed from Theorem 2 and Remark 1 it is immediate that  $m_*(X; Y) \geq \rho_m^2(X; Y)$ . We will prove later in the paper that the inequality is strict in general by providing an explicit example. The strictness of the inequality  $r_\infty(X; Y) \geq \rho_m^2(X; Y)$  in general is known from [1, Theorem 9b]. The strictness of the inequality  $m_*(X; Y) \geq \rho_m^2(X; Y)$  would also alternately follow from our Theorem 2.*

In this paper we will also provide alternate geometric characterizations of both  $\rho_m^2(X; Y)$  and  $r_\infty(X; Y)$ . Fix a channel  $\mu_{Y|X}(y|x)$ , fix  $\lambda \in [0, 1]$ , and consider the function<sup>2</sup> of the probability distribution of  $X$  denoted by  $t_\lambda(X)$  which is defined by

$$t_\lambda(X) := H(Y) - \lambda H(X).$$

<sup>2</sup>We abuse notation when we write  $t_\lambda(X)$ . We really wish to think of  $t_\lambda$  as a function of the probability distribution of  $X$ .

We will show in Theorem 3 that  $\rho_m^2(X; Y)$  is the smallest  $\lambda$  such that  $t_\lambda(X)$  has a positive semidefinite Hessian at  $\mu(x)$  and  $r_\infty(X; Y)$  is the smallest  $\lambda$  such that  $t_\lambda(X)$  matches its lower convex envelope, denoted by  $\mathcal{K}[t_\lambda](X)$ , at  $\mu_X(x)$ .

### C. Organization of the paper

In Section II, we will prove Theorem 2. In Section III we will establish the alternate geometric characterizations for  $\rho_m^2(X; Y)$  and  $r_\infty(X; Y)$ . The paper we uploaded on arXiv [15] presents the key results from a slightly different perspective.

## II. PROOF OF THEOREM 2

*Remark:* The proof below is stitched together using a judicious borrowing of arguments from [1], [14], and standard techniques. In the authors' opinion, the three way equivalence argument elucidates the proof of the equivalence  $r_\infty(X; Y) = d_*(X; Y)$  (in [1]). Further this proof idea allows for a natural generalization and provides an alternate (new) characterization of  $r_p(X; Y)$ ,  $p \geq 1$  [18].

If  $X$  and  $Y$  are independent, then it is easy to see that  $r_p(X; Y) = \frac{1}{p} \forall p \geq 1$ ; hence  $r_\infty(X; Y) = 0$ . It is also easy to see that  $d_*(X; Y) = m_*(X; Y) = 0$ . The theorem clearly holds in this case.

We will assume then that  $X$  and  $Y$  are not independent. By choosing  $f(x) = 1_{x \in A} - P(X \in A)$ ,  $g(y) = 1_{y \in B} - P(Y \in B)$  in (1) for appropriate sets  $A, B$ , we can obtain  $\rho_m(X; Y) > 0$ . From (3), we get  $r_\infty(X; Y) > 0$ .

The proof will follow from the following sequence of implications that we will establish.

- (a)  $r_\infty(X; Y) \leq d_*(X; Y)$ ,
- (b)  $d_*(X; Y) \leq m_*(X; Y)$ ,
- (c)  $m_*(X; Y) \leq r_\infty(X; Y)$ .

*Proof of (a):* Writing  $r_p, r_\infty$  for  $r_p(X; Y), r_\infty(X; Y)$  respectively, from the definition we have the following inequality:

$$\|E(g(Y)|X)\|_p \leq \|g(Y)\|_{r_p p},$$

for all  $g(Y) \geq 0$ .

Define  $h(y) := g(y)^{r_p p}$ , and note that

$$\|g(Y)\|_{r_p p}^{r_p p} = E(h(Y)).$$

We also obtain that for all  $p \geq 1$

$$\|E(g(Y)|X)\|_p^{r_p p} = \left( \sum_x \mu(x) \left( \sum_y \mu(y|x) h(y)^{\frac{1}{r_p p}} \right)^p \right)^{r_p}.$$

Using the well-known limit  $\lim_{r \downarrow 0} \|W\|_r = \exp(E \log |W|)$ , we get that as  $p \rightarrow \infty$ ,

$$\left( \sum_y \mu(y|x) h(y)^{\frac{1}{r_p p}} \right)^p \rightarrow \prod_y h(y)^{\frac{\mu(y|x)}{r_\infty}}.$$

Thus we have for all  $h(Y) > 0$

$$\left( \sum_x \mu(x) \prod_y h(y)^{\frac{\mu(y|x)}{r_\infty}} \right)^{r_\infty} \leq E(h(Y)). \quad (6)$$

From the definition of  $r_\infty(X; Y)$  and the continuity and strict monotonicity of  $\|\cdot\|_r$  in  $r$  for non-constant random variables, it is immediate that it is the smallest such number for which the above inequality holds for all  $h(Y) > 0$ .

Therefore for any  $\epsilon > 0$  there exists an  $h_\epsilon(Y)$  such that

$$\left( \sum_x \mu(x) \prod_y h_\epsilon(y) \frac{\mu(y|x)}{r_\infty - \epsilon} \right)^{r_\infty - \epsilon} > \mathbb{E}(h_\epsilon(Y)).$$

Further w.l.o.g. assume that

$$\mathbb{E}(h_\epsilon(Y)) = 1.$$

Define a probability distribution (see [1, (5.11)])

$$\nu(x) = C \mu(x) \prod_y h_\epsilon(y) \frac{\mu(y|x)}{r_\infty - \epsilon},$$

where  $C < 1$  is a normalizing constant. Now, note that

$$\nu(x) \log \frac{\nu(x)}{C \mu(x)} = \nu(x) \sum_y \frac{\mu(y|x)}{r_\infty - \epsilon} \log h_\epsilon(y)$$

$$\implies \sum_x \nu(x) \log \frac{\nu(x)}{\mu(x)} = \sum_y \frac{\nu^\mu(y)}{r_\infty - \epsilon} \log h_\epsilon(y) + \log C.$$

Finally observe that since  $C < 1$

$$\begin{aligned} & \sum_x \nu(x) \log \frac{\nu(x)}{\mu(x)} \\ & < \sum_y \frac{\nu^\mu(y)}{r_\infty - \epsilon} \log h_\epsilon(y) \\ & = \sum_y \frac{\nu^\mu(y)}{r_\infty - \epsilon} \log \frac{\nu^\mu(y)}{\mu(y)} + \sum_y \frac{\nu^\mu(y)}{r_\infty - \epsilon} \log \frac{\mu(y) h_\epsilon(y)}{\nu^\mu(y)} \\ & \leq \sum_y \frac{\nu^\mu(y)}{r_\infty - \epsilon} \log \frac{\nu^\mu(y)}{\mu(y)} + \frac{1}{r_\infty - \epsilon} \log \left( \sum_y \mu(y) h_\epsilon(y) \right) \\ & = \sum_y \frac{\nu^\mu(y)}{r_\infty - \epsilon} \log \frac{\nu^\mu(y)}{\mu(y)}, \end{aligned}$$

where the last equality follows since  $\mathbb{E}(h_\epsilon(Y)) = 1$ .

Thus  $r_\infty(X; Y) - \epsilon < \frac{D(\nu^\mu(y) \|\mu(y))}{D(\nu^\mu(y) \|\mu(x))} \leq d_*(X; Y)$ . Since  $\epsilon > 0$  is arbitrary we are done.

*Proof of (b):* Let  $\mathcal{U}_\epsilon := \{1, 2\}$ . Fix a sufficiently small  $\epsilon > 0$  and define  $U_\epsilon$  satisfying  $U_\epsilon - X - Y$  by

- $P(U_\epsilon = 1) = \epsilon, P(X = x | U_\epsilon = 1) = \nu_\delta(x),$
- $P(U_\epsilon = 2) = 1 - \epsilon, P(X = x | U_\epsilon = 2) = \mu(x) + \frac{\epsilon}{1-\epsilon}(\mu(x) - \nu_\delta(x)) = \frac{1}{1-\epsilon}\mu(x) - \frac{\epsilon}{1-\epsilon}\nu_\delta(x),$

where  $\nu_\delta(x) \neq \mu(x)$  is a probability distribution satisfying  $\frac{D(\nu_\delta^\mu(y) \|\mu(y))}{D(\mu_\delta(x) \|\mu(x))} > d_*(X; Y) - \delta > 0$ . For sufficiently small  $\epsilon > 0$ , we have that  $\frac{1}{1-\epsilon}\mu(x) - \frac{\epsilon}{1-\epsilon}\nu_\delta(x)$  is a probability distribution (as  $\mu(x)$  was assumed to have full support). Note that

$$\begin{aligned} & P(U_\epsilon = 1)P(X = x | U_\epsilon = 1) + \\ & P(U_\epsilon = 2)P(X = x | U_\epsilon = 2) = \mu(x) \quad \forall x \in \mathcal{X}, \end{aligned}$$

so that this specified chain  $U_\epsilon - X - Y$  has the correct marginal distribution for  $(X, Y)$ .

For any  $0 < \theta < d^*(X; Y) - \delta$  define the function

$$g(\epsilon) := I(U_\epsilon; Y) - \theta I(U_\epsilon; X).$$

We have

$$\begin{aligned} \frac{dg(\epsilon)}{d\epsilon} &= -\frac{d}{d\epsilon} \left( \epsilon H(\nu_\delta^\mu(y)) + (1-\epsilon)H \left( \frac{1}{1-\epsilon}\mu(y) - \frac{\epsilon}{1-\epsilon}\nu_\delta^\mu(y) \right) \right) \\ &+ \theta \frac{d}{d\epsilon} \left( \epsilon H(\nu_\delta(x)) + (1-\epsilon)H \left( \frac{1}{1-\epsilon}\mu(x) - \frac{\epsilon}{1-\epsilon}\nu_\delta(x) \right) \right) \\ &= -H(\nu_\delta^\mu(y)) + H \left( \frac{\mu(y) - \epsilon \nu_\delta^\mu(y)}{1-\epsilon} \right) + \theta H(\nu_\delta(x)) \\ &- \theta H \left( \frac{\mu(x) - \epsilon \nu_\delta(x)}{1-\epsilon} \right) - \sum_y \frac{\nu_\delta^\mu(y) - \mu(y)}{1-\epsilon} \log \left( \frac{\mu(y) - \epsilon \nu_\delta^\mu(y)}{1-\epsilon} \right) \\ &+ \theta \sum_x \frac{\nu_\delta(x) - \mu(x)}{1-\epsilon} \log \left( \frac{\mu(x) - \epsilon \nu_\delta(x)}{1-\epsilon} \right). \end{aligned}$$

Thus

$$\left. \frac{dg(\epsilon)}{d\epsilon} \right|_{\epsilon=0} = D(\nu_\delta^\mu(y) \|\mu(y)) - \theta D(\nu_\delta(x) \|\mu(x)) > 0,$$

where the last inequality is because  $0 < \theta < d^*(X; Y) - \delta$  and  $\frac{D(\nu_\delta^\mu(y) \|\mu(y))}{D(\nu_\delta(x) \|\mu(x))} > d_*(X; Y) - \delta$ . Since  $g(0) = 0$  this implies that for some  $\epsilon' > 0$  we have  $I(U_{\epsilon'}; Y) - \theta I(U_{\epsilon'}; X) > 0$  or that

$$m_*(X; Y) = \sup_{U: U-X-Y, I(U; Y) > 0} \frac{I(U; Y)}{I(U; X)} \geq \frac{I(U_{\epsilon'}; Y)}{I(U_{\epsilon'}; X)} > \theta.$$

Since the above holds for all  $\theta < d_*(X; Y) - \delta$  we have

$$m_*(X; Y) \geq d_*(X; Y) - \delta.$$

Finally, since  $\delta > 0$  is arbitrary, we let  $\delta \rightarrow 0$ , and we are done.

*Proof of (c):* This part uses standard typicality arguments in information theory and our definition of (and notation for)  $\epsilon$ -typical sets are borrowed from [16].

For any  $U \rightarrow X \rightarrow Y$  let  $(U^n, X^n, Y^n) \sim \prod_i \mu(u_i, x_i, y_i)$ . Pick a single  $u^n \in \mathcal{T}_\epsilon^{(n)}(U)$ . For some  $\epsilon_1 > \epsilon$  let  $\mathcal{A}_n = \{x^n : (u^n, x^n) \in \mathcal{T}_{\epsilon_1}^{(n)}(U, X)\}$  and for  $\epsilon_2 > \epsilon_1$  let  $\mathcal{B}_n = \{y^n : (u^n, y^n) \in \mathcal{T}_{\epsilon_2}^{(n)}(U, Y)\}$ . Note that

$$\begin{aligned} & P(X^n \in \mathcal{A}_n, Y^n \in \mathcal{B}_n) \\ &= \mathbb{E}[1_{X^n \in \mathcal{A}_n} \mathbb{E}(1_{Y^n \in \mathcal{B}_n} | X^n)] \\ &\leq \mathbb{E} \|1_{X^n \in \mathcal{A}_n}\|_{\frac{p}{p-1}} \mathbb{E} \|1_{Y^n \in \mathcal{B}_n} | X^n\|_p \\ &\leq \mathbb{E} \|1_{X^n \in \mathcal{A}_n}\|_{\frac{p}{p-1}} \mathbb{E} \|1_{Y^n \in \mathcal{B}_n}\|_{r_p p} \\ &= P(X^n \in \mathcal{A}_n)^{1-\frac{1}{p}} P(Y^n \in \mathcal{B}_n)^{\frac{1}{r_p p}}, \end{aligned}$$

where we write  $r_p$  to denote  $r_p(X; Y)$  for convenience. The first inequality follows from Hölder's inequality and the second inequality from the definition and tensorization property of  $r_p(X; Y)$ .

Standard calculations tell us that  $\frac{1}{n} \log_2 P(X^n \in \mathcal{A}_n) \rightarrow -I(U; X)$  and  $\frac{1}{n} \log_2 P(Y^n \in \mathcal{B}_n) \rightarrow -I(U; Y)$  as  $n \rightarrow \infty$ . From the law of large numbers we know that  $P(Y^n \in \mathcal{B}_n | X^n \in \mathcal{A}_n) \rightarrow 1$  as  $n \rightarrow \infty$ . Therefore, taking the

logarithm on both sides, dividing by  $n$  and letting  $n \rightarrow \infty$  we obtain that

$$-I(U; X) \leq -\left(1 - \frac{1}{p}\right) I(U; X) - \frac{1}{r_p p} I(U; Y).$$

Rearranging we obtain  $r_p(X; Y) \geq \frac{I(U; Y)}{I(U; X)}$  and since this is true for any  $U$ , by taking the supremum on the right hand side we obtain  $r_p(X; Y) \geq m_*(X; Y)$ . Since the right hand side does not depend on  $p$ , we take  $p \rightarrow \infty$  to obtain  $r_\infty(X; Y) \geq m_*(X; Y)$ . This completes the proof of Theorem 2.

### III. A GEOMETRIC CHARACTERIZATION OF $\rho_m^2(X; Y)$ AND $m_*(X; Y)$

Let  $\mu(y|x)$  be the channel transition probability from  $X$  to  $Y$  induced by a joint distribution  $\mu(x, y)$ . For a fixed channel  $\mu(y|x)$ , consider a function of the input distribution  $\nu(x)$ ,

$$t_\lambda(X) := H(Y) - \lambda H(X),$$

where  $\lambda$  is a constant in  $[0, 1]$ . Observe that the function is concave when  $\lambda = 0$  and convex when  $\lambda = 1$ .<sup>3</sup>

We write  $\mathcal{K}[t_\lambda](X)$  for the lower convex envelope of  $t_\lambda(X)$ .

**Proposition 1.** *If  $\mathcal{K}[t_\lambda](X) = t_\lambda(X)$  at  $\nu(x)$  for some  $\lambda$  then  $\mathcal{K}[t_{\lambda_1}](X) = t_{\lambda_1}(X)$  at  $\nu(x)$  for all  $\lambda_1 \geq \lambda$ .*

*Proof.* If  $\mathcal{K}[t_\lambda](X) = t_\lambda(X)$  at  $\nu(x)$  for some  $\lambda$ , then note that for any  $\lambda_1 \geq \lambda$

$$\begin{aligned} t_{\lambda_1}(X) &= t_\lambda(X) - (\lambda_1 - \lambda)H(X) \\ \implies \mathcal{K}[t_{\lambda_1}](X) &\geq \mathcal{K}[t_\lambda](X) - (\lambda_1 - \lambda)H(X). \end{aligned}$$

Here the inequality comes since  $\mathcal{K}[f+g] \geq \mathcal{K}[f] + \mathcal{K}[g]$ ; note that  $-(\lambda_1 - \lambda)H(X)$  is convex. Therefore at  $\nu(x)$  we will have that

$$\begin{aligned} t_{\lambda_1}(X) &\geq \mathcal{K}[t_{\lambda_1}](X) \geq \mathcal{K}[t_\lambda](X) - (\lambda_1 - \lambda)H(X) \\ &= t_\lambda(X) - (\lambda_1 - \lambda)H(X) = t_{\lambda_1}(X), \end{aligned}$$

establishing the proposition.  $\square$

The following theorem gives a geometric interpretation of  $\rho_m^2(X; Y)$  and  $m_*(X; Y)$  in terms of the behaviour of the function  $t_\lambda(X)$ .

**Theorem 3.** *Let  $(X, Y) \sim \mu(x, y)$ . The following statements hold:*

- 1)  $\rho_m^2(X; Y)$  is the minimum value of  $\lambda$  such that the function  $t_\lambda(X)$  has a positive semidefinite Hessian at  $\mu(x)$ .
- 2)  $m_*(X; Y)$  is the minimum value of  $\lambda$  such that the function  $t_\lambda(X)$  touches its lower convex envelope at  $\mu(x)$ , i.e. such that  $\mathcal{K}[t_\lambda](X) = t_\lambda(X)$  at  $\mu(x)$ .

*Proof of 1):* The claim is straightforward when  $X, Y$  are independent. When  $X, Y$  are not independent, Rényi's characterization of the maximal correlation [19] states that

$$\rho_m^2(X; Y) = \sup_{f(X): \mathbb{E} f(X)=0, \mathbb{E}[f^2(X)]=1} \mathbb{E}[\mathbb{E}[f(X)|Y]^2].$$

<sup>3</sup>This convexity at  $\lambda = 1$  follows from the fact that for any  $U - X - Y$  we have  $I(U; X) \geq I(U; Y)$  or equivalently  $H(Y) - H(X) \leq H(Y|U) - H(X|U)$ .

Take an arbitrary multiplicative perturbation of the form  $\mu_\epsilon(x) = \mu(x)(1 + \epsilon f(x))$ . For  $\mu_\epsilon$  to stay a valid perturbation we need  $\mathbb{E}[f(X)] = 0$ . Furthermore we can normalize  $f$  by assuming that  $\mathbb{E}[f^2(X)] = 1$ . The second derivative in  $\epsilon$  of  $H(Y) - \lambda H(X)$  is equal to [7]

$$-\mathbb{E}[\mathbb{E}[f(X)|Y]^2] + \lambda \mathbb{E}[f^2(X)] = -\mathbb{E}[\mathbb{E}[f(X)|Y]^2] + \lambda,$$

which is non-negative as long as  $\lambda \geq \mathbb{E}[\mathbb{E}[f(X)|Y]^2]$ . Thus the minimum value  $\lambda^*$  such that the second derivative is non-negative for all local perturbations is

$$\lambda^* = \sup_{f(X): \mathbb{E} f(X)=0, \mathbb{E}[f^2(X)]=1} \mathbb{E}[\mathbb{E}[f(X)|Y]^2] = \rho_m^2(X; Y). \quad \blacksquare$$

*Proof of 2):* Consider the minimum value of  $\lambda$ , say  $\lambda^\dagger$ , such that the function  $t_\lambda(X)$  touches its lower convex envelope at  $\mu(x)$ . Thus, we are looking for the minimum  $\lambda$  such that for  $(X, Y) \sim \mu(x, y)$  we have<sup>4</sup>

$$H(Y) - \lambda H(X) \leq H(Y|U) - \lambda H(X|U), \quad \forall U : U - X - Y.$$

Equivalently we require the minimum  $\lambda$  such that,

$$\lambda \geq \frac{I(U; Y)}{I(U; X)}, \quad \forall U : U - X - Y \text{ with } I(U; X) > 0.$$

Thus,

$$\lambda^\dagger = \sup_{U: U-X-Y, I(U; X) > 0} \frac{I(U; Y)}{I(U; X)} = m_*(X; Y). \quad \blacksquare$$

**Remark 5.** *Since  $t_\lambda(X) = \mathcal{K}[t_\lambda](X)$  at  $\mu(x)$  implies that the Hessian of  $t_\lambda(X)$  at  $\mu(x)$  is positive semidefinite, we have*

$$m_*(X; Y) \geq \rho_m^2(X; Y).$$

**A. Counterexample to the Erkip-Cover data processing inequality**

In [3, Theorem 8], Erkip and Cover claimed that

$$I(U; Y) \leq \rho_m^2(X; Y) I(U; X)$$

holds whenever  $U - X - Y$  form a Markov chain. Furthermore they claimed that,  $\rho_m^2(X; Y)$  is the minimum such constant, i.e.

$$m_*(X; Y) = \rho_m^2(X; Y). \quad (7)$$

We will first provide a counterexample to these claims and then point out a gap in their argument.

1) *Counterexample to (7):* Let  $X$  be a binary random variable with  $p(X = 0) = \frac{1}{2}$ . Define  $p(x, y)$  by passing  $X$  through the asymmetric erasure channel given in Fig. 1. When either  $X$  or  $Y$  is binary then it is known [20] that

$$\rho_m^2(X; Y) = \left[ \sum_{x,y} \frac{p(x,y)^2}{p(x)p(y)} \right] - 1.$$

We then have  $\rho_m^2(X; Y) = 0.6$ . Suppose we construct  $U$  satisfying  $U - X - Y$  such that  $U|\{X = 0\} \sim \text{Ber}(0.1)$ ,  $U|\{X = 1\} \sim \text{Ber}(0.4)$ . Then  $I(U; Y) = 0.055770\dots$  and  $I(U; X) =$

<sup>4</sup>Note that if  $U$  is independent of  $X$ , i.e.  $I(U; X) = 0$  then the above inequality is always true.

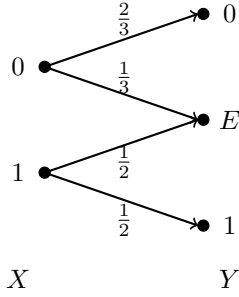


Fig. 1. An asymmetric erasure channel.

0.09130..., so that  $\frac{I(Y;U)}{I(X;U)} = 0.6108... > 0.6 = \rho_m^2(X;Y)$ , and this contradicts (7). It can be shown in a reasonably straightforward manner, using our characterization in Theorem 3, that  $m_*(X;Y) = \frac{1}{2} \log_2\left(\frac{12}{5}\right) = 0.631517...$  for this pair of random variables  $(X, Y)$ .

One can show that if a measure  $\mu(x, y)$  is drawn uniformly at random from the set of all probability measures on pairs of binary random variables, then with probability one we have  $m_*(X;Y) > \rho_m^2(X;Y)$ .

The error of the Erkip-Cover proof lies in their use of a Taylor’s series expansion. Consider the expansion in the left column of page 1037 of their paper [3], where they use their equation (16) to expand around  $p(\tilde{v})$ . It is possible that  $p(\tilde{v})$  is zero for some  $\tilde{v}$  and this causes an error as the derivative in this direction is infinity and the Taylor’s series expansion is no longer valid. As our counterexample shows, this seems to be a significant but subtle error that cannot be worked around.

Some of the works that use this incorrect result of [3], such as [21], are affected by this error. A claim similar to that of [3], which appears in [9], is also false.<sup>5</sup>

#### IV. CONCLUSION

In this paper we showed the equivalence between the optimal constant in the data-processing inequality and a hypercontractivity parameter connecting random variables  $X$  and  $Y$ . This corrects an incorrect claim due to Erkip and Cover [3]. We also presented a new geometric characterization of the maximal correlation and of this hypercontractivity parameter.

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<sup>5</sup>This paper studies the ratio  $\frac{I(U;Y)}{I(U;X)}$  when  $I(U;X)$  is very small. However, as pointed out in [3], the supremum of  $\frac{I(U;Y)}{I(U;X)}$  occurs when  $I(U;X) \rightarrow 0$ . So the problem studied by [9] is the same as that of [3].

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